

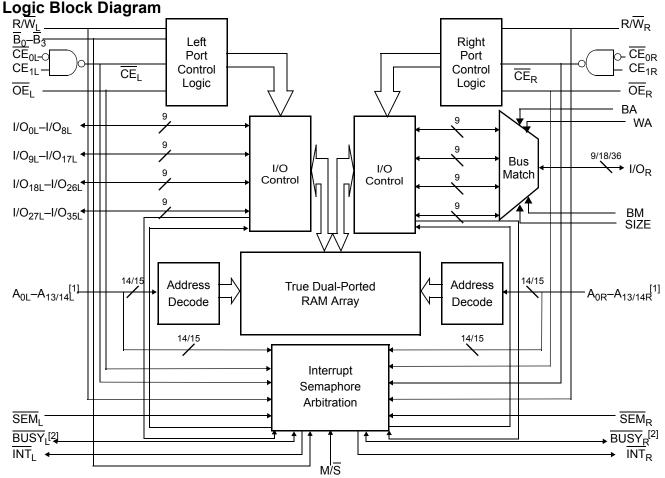


3.3 V 16K/32K x 36 FLEx36™ Asynchronous Dual-Port Static **RAM**

Features

- True dual-ported memory cells that allow simultaneous access of the same memory location
- 16K x 36 organization (CY7C056V)
- 32K x 36 organization (CY7C057V)
- 0.25-micron Complimentary metal oxide semiconductor (CMOS) for optimum speed/power
- High-speed access: 12/15 ns
- Low operating power
 - ☐ Active: I_{CC} = 250 mA (typical)
 - Standby: I_{SB3} = 10 μA (typical)
- Fully asynchronous operation
- Automatic power-down
- Expandable data bus to 72 bits or more using Master/Slave Chip Select when using more than one device

- On-chip arbitration logic
- Semaphores included to permit software handshaking between ports
- INT flag for port-to-port communication
- Byte select on left port
- Bus matching on right port
- Depth expansion via dual chip enables
- Pin select for Master or Slave
- Commercial and Industrial temperature ranges
- Available in 144-Pin Thin quad plastic flatpack (TQFP) or 172-Ball ball grid array (BGA)
- Pb-free packages available
- Compact packages:
 - □ 144-Pin TQFP (20 x 20 x 1.4 mm)
 - □ 172-Ball BGA (1.0-mm pitch) (15 x 15 x.51 mm)



- A₀-A₁₃ for 16K; A₀-A₁₄ for 32K devices.
 BUSY is an output in Master mode and an input in Slave mode.

[+] Feedback



Contents

Pin Configurations	4
Pin Configurations	5
Selection Guide	
Pin Definitions	6
Maximum Ratings	7
Operating Range	7
Electrical Characteristics	8
Capacitance	8
AC Test Load and Waveforms	9
Switching Characteristics	. 10
Data Retention Mode	. 11
Timing	
Read Cycle No. 1 (Either Port Address Access)	
Read Cycle No. 2 (Either Port CE/OE Access)	
Read Cycle No. 3 (Either Port)	
Switching Waveforms	
Write Cycle No. 1: R/W Controlled Timing	
Write Cycle No. 2: CE Controlled Timing	
Semaphore Read After Write Timing, Either Side	
Timing Diagram of Semaphore Contention	
Timing Diagram of Write with BUSY (M/S = HIGH)	
Write Timing with Busy Input $(M/\overline{S} = LOW)$	
Busy Timing Diagram No. 1 (CE Arbitration)	. 16
Busy Timing Diagram No. 2 (Address Arbitration)	. 16
Architecture	18

Functional Description	
Write Operation	18
Read Operation	18
Interrupts	18
Busy	18
Master/Slave	18
Semaphore Operation	18
Right Port Configuration	20
Right Port Operation	20
Left Port Operation	20
Bus Match Operation	20
Long-Word (36-bit) Operation	21
Word (18-bit) Operation	
Byte (9-bit) Operation	21
Ordering Information	22
Ordering Code Definition	22
Package Diagrams	23
Acronyms	24
Document Conventions	24
Units of Measure	24
Sales, Solutions, and Legal Information	26
Worldwide Sales and Design Support	
Products	
PSoC Solutions	26



Functional Description

The CY7C056V and CY7C057V are low-power CMOS 16K and 32K x 36 dual-port static RAMs. Various arbitration schemes are included on the devices to handle situations when multiple processors access the same piece of data. Two ports are provided, permitting independent, asynchronous access for reads and writes to any location in memory. The devices can be utilized as standalone 36-bit dual-port static RAMs or multiple devices can be combined in order to function as a 72-bit or wider master/slave dual-port static RAM. An M/S pin is provided for implementing 72-bit or wider memory applications without the need for separate master and slave devices or additional discrete logic. Application areas include interprocessor/multiprocessor designs, communications status buffering, and dual-port video/graphics memory.

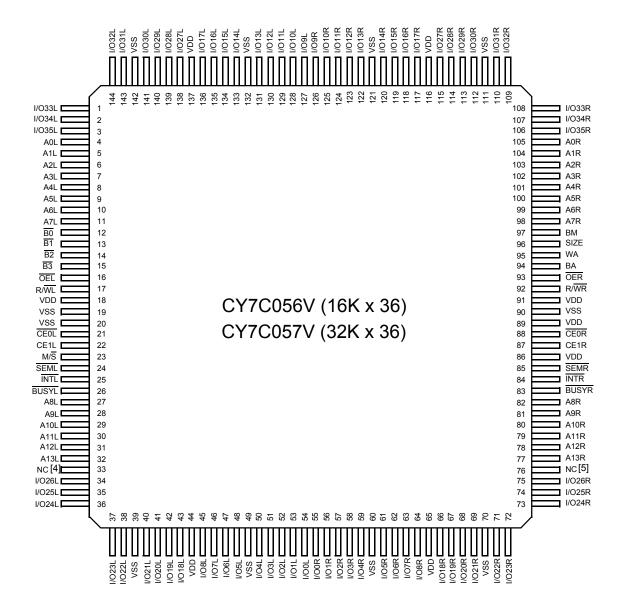
Each port has independent control pins: Chip Enable $(\overline{CE})^{[3]}$, Read or Write Enable (R/W), and Output Enable (\overline{OE}) . Two flags are provided on each port (BUSY) and \overline{INT} . BUSY signals that the port is trying to access the same location currently being accessed by the other port. The Interrupt Flag (\overline{INT}) permits communication between ports or systems by means of a mail box. The semaphores are used to pass a flag, or token, from one port to the other to indicate that a shared resource is in use. The semaphore logic is comprised of eight shared latches. Only one side can control the latch (semaphore) at any time. Control of a semaphore indicates that a shared resource is in use. An automatic Power-down feature is controlled independently on each port by Chip Select (\overline{CE}_0) and \overline{CE}_1 pins.

The CY7C056V and CY7C057V are available in 144-Pin Thin quad plastic flatpack (TQFP) and 172-Ball ball grid array (BGA) packages.



Pin Configurations

Figure 1. 144-Pin Thin Quad Flatpack (TQFP) Top View



- 4. This pin is A14L for CY7C057V.
- 5. This pin is A14R for CY7C057V



Pin Configurations (continued)

Figure 2. 172-Ball Ball Grid Array (BGA)
Top View

1 2 6 7 3 4 5 8 9 10 11 12 13 14 I/O32L I/O30L I/O13L VDD I/O11L I/011R VDD I/O13R VSS NC I/O30R I/032R I/O14L I/O12L I/O14R I/017R I/033R В A0L I/O33L I/O29 I/O17L I/O9L I/09R I/012R I/O29R A0R I/O27L I/O10L I/O10R I/O15R I/027R NC NC I/O31L NC I/O15L NC I/031R A1R С A1L A2L A3L I/O35L I/O34L I/O28L I/O16L VSS I/O16R I/028R I/O34R I/O35R A3R A2R D Ε A4L A5L B0L A4R B1L F VDD A6L A7L NC NC SIZE A7R A6R VDD OEL B2L B3L CE0L CE0R WA OER ВА G R/WL A8L CE1L CE1R A8R R/WR VSS Н M/S A9L A10L VSS NC NC VDD VDD A10R A9R A11L A12L NC SEML NC NC NC NC SEMR NC A12R A11R BUSYL INTL I/O26L I/O25L I/O19L VSS VSS I/O19R I/O25R I/O26R INTR BUSYR A13L A13R L M I/O22L I/O18L NC I/O7L I/O2L I/07R I/O18R I/022R NC^[5] NC I/O20L I/O24L I/O8L I/O6L I/O5L I/O0L I/O0R I/O3R I/O5R I/08R I/O20R I/O24R I/O3L I/06R I/O23L I/O21L VDD I/021R I/023R NC VSS I/O4L VDD I/O1L I/O1R I/O4R VSS NC

Selection Guide

	CY7C056V CY7C057V -12	CY7C056V CY7C057V -15	Unit
Maximum access time	12	15	ns
Typical operating current	250	240	mA
Typical standby current for I _{SB1} (Both ports TTL level)	55	50	mA
Typical standby current for I _{SB3} (Both ports CMOS level)	10 μΑ	10 μΑ	μΑ



Pin Definitions

Left Port	Right Port	Description
A _{0L} -A _{13/14L}	A _{0R} -A _{13/14R}	Address (A ₀ -A ₁₃ for 16K; A ₀ -A ₁₄ for 32K devices)
SEM _L	SEM _R	Semaphore Enable
CE _{0L} , CE _{1L}	CE _{0R} , CE _{1R}	Chip Enable (\overline{CE} is LOW when $\overline{CE}_0 \le V_{IL}$ and $CE_1 \ge V_{IH}$)
ĪNT _L	INT _R	Interrupt flag
BUSY _L	BUSY _R	Busy flag
I/O _{0L} –I/O _{35L}	I/O _{0R} –I/O _{35R}	Data bus input/output
ŌĒL	OE _R	Output Enable
R/\overline{W}_L	R/W _R	Read/Write Enable
$\overline{B}_0 - \overline{B}_3$		Byte select inputs. Asserting these signals enables read and write operations to the corresponding bytes of the memory array.
	BM, SIZE	See bus matching for details.
	WA, BA	See bus matching for details.
M/S		Master or Slave select
V _{SS}		Ground
V_{DD}		Power



Maximum Ratings^[6]

Output current into outputs (LOW)	20 mA
Static discharge voltage	>2001 V
Latch-up current	>200 mA

Operating Range

Range	Ambient Temperature	V _{DD}
Commercial	0 °C to +70 °C	3.3 V ± 165 mV
Industrial	–40 °C to +85 °C	3.3 V ± 165 mV

Notes

6. The voltage on any input or I/O pin can not exceed the power pin during power-up.

7. Pulse width < 20 ns.



Electrical Characteristics Over the Operating Range^[8]

					056V 057V				
Parameter	rameter Description		Description -12			-15			
			Min	Тур	Max	Min	Тур	Max	Unit
V _{OH}	Output HIGH voltage (V _{DD} = Min., I _{OH} = -4.0 mA)				_	2.4		-	V
V _{OL}	Output LOW voltage (V _{DD} = Min., I _{OL} = +4.0 mA)		_		0.4	-	_	0.4	V
V _{IH}	Input HIGH voltage		2.0		_	2.0		_	V
V _{IL}	Input LOW voltage		_		0.8	_		8.0	V
I _{OZ}	Output leakage current				10	-10		10	μА
I _{CC}	Operating current (V _{DD} = Max.,	Commercial		250	385		240	360	mA
	I _{OUT} = 0 mA) output disabled	Industrial		_			265	385	mA
I _{SB1}	Standby current (Both ports TTL level	Commercial		55	75		50	70	mA
	and deselected) f = f _{MAX}	Industrial		_			65	95	mA
I _{SB2}	Standby current (One port TTL	Commercial		180	240		175	230	mA
	level and deselected) f = f _{MAX}	Industrial	-		-		190	255	mA
I _{SB3}	Standby current (Both ports CMOS	Commercial		0.01	1		0.01	1	mA
	level and deselected) f =0			-	-		0.01	1	mA
I _{SB4}	Standby current (One Port CMOS	Commercial		160	210		155	200	mA
	level and deselected) f = f _{MAX} ^[9]	Industrial		-	-		170	215	mA

Capacitance^[10]

Parameter	Description	Test Conditions	Max	Unit
C _{IN}	Input capacitance	$T_A = 25 ^{\circ}\text{C}, f = 1 \text{MHz},$	10	pF
C _{OUT}	Output capacitance	V _{DD} = 3.3 V	10	pF

^{8.} Deselection for a port occurs if $\overline{\text{CE}}_0$ is HIGH or if CE_1 is LOW.

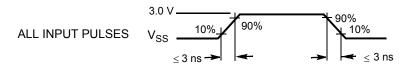
9. $f_{\text{MAX}} = 1/t_{\text{RC}} = \text{All inputs cycling at } f = 1/t_{\text{RC}}$ (except Output Enable). f = 0 means no address or control lines change. This applies only to inputs at CMOS level standby I_{SB3} .

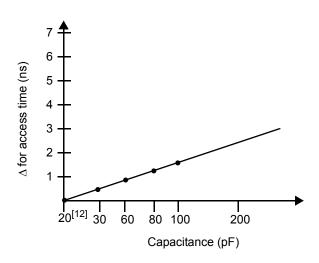
10. Tested initially and after any design or process changes that may affect these parameters.



AC Test Load and Waveforms







(b) Load Derating Curve

^{11.} External AC Test Load Capacitance = 10 pF.
12. (Internal I/O pad Capacitance = 10 pF) + AC Test Load.



Switching Characteristics Over the Operating Range^[13]

_						
Parameter	Description	-	12		15	Unit
		Min	Max	Min	Max	
Read Cycle		•				
t _{RC}	Read cycle time	12	_	15	-	ns
t _{AA}	Address to data valid	_	12	-	15	ns
t _{OHA}	Output hold from address change	3	_	3	_	ns
t _{ACE} ^[14, 15]	CE LOW to data valid	_	12	_	15	ns
tnoe	OE LOW to data valid	_	8	_	10	ns
t _{LZOE} [14, 16, 17, 18]	OE Low to low Z	0	_	0	_	ns
t _{HZOE} [14, 16, 17, 18]	OE HIGH to High Z	_	10	-	10	ns
t _{LZCE} [14, 13, 17, 18]	CE LOW to Low Z	3	_	3	_	ns
t _{HZCE} [14, 16, 17, 18]		_	10	-	10	ns
t _{LZBE}	Byte Enable to Low Z	3	_	3	_	ns
t _{HZBE}	Byte Enable to High Z	_	10	-	10	ns
t _{PU} ^[14, 18]	CE LOW to power-up	0	_	0	_	ns
t _{PD} ^[14, 18]	CE HIGH to power-down	_	12	-	15	ns
t _{ABE} ^[15]	Byte Enable access time	_	12	-	15	ns
Write Cycle	1	-	ı	l		
t _{WC}	Write cycle time	12	_	15	_	ns
t _{SCE} ^[14, 15]	CE LOW to write end	10	_	12	_	ns
t _{AW}	Address valid to write end	10	_	12	_	ns
t _{HA}	Address hold from write end	0	_	0	_	ns
t _{SA} ^[15]	Address set-up to write start	0	_	0	_	ns
t _{PWE}	Write pulse width	10	_	12	_	ns
t _{SD}	Data set-up to write end	10	_	10	_	ns
t _{HD}	Data hold from write end	0	_	0	_	ns
t _{HZWE} ^[17, 18]	R/W LOW to High Z	_	10	-	_	ns
t _{LZWE} [17, 18]	R/W HIGH to Low Z	3	-	3	_	ns
t _{WDD} ^[19]	Write pulse to data delay	_	25	-	_	ns
t _{DDD} ^[19]	Write data valid to read data valid	_	20	_	25	ns
Busy Timing ^[20]		1	ı			
t _{BLA}	BUSY LOW from address match	_	12	_	15	ns
t _{BHA}	BUSY HIGH from address mismatch	_	12	-	15	ns
t _{BLC}	BUSY LOW from CE LOW	_	12	-	15	ns
Notes	1	_1	1	1	11	

- 13. Test conditions assume signal transition time of 3 ns or less, timing reference levels of 1.5 V, input pulse levels of 0 to 3.0 V, and output loading of the specified I_{OI}/I_{OH} and 10-pF load capacitance.

 14. CE is LOW when CE₀ ≤ V_{IL} and CE₁ ≥ V_{IH}

 15. To access RAM, CE = L and SEM = H. To access semaphore, CE = H and SEM = L. Either condition must be valid for the entire t_{SCE} time.

- 16. At any given temperature and voltage condition for any given device, t_{HZCE} is less than t_{LZCE} and t_{HZOE} is less than t_{LZCE}.

 17. Test conditions used are Load 2.
- 18. This parameter is guaranteed by design, but it is not production tested. For information on port-to-port delay through RAM cells from writing port to reading port, refer to Read Timing with Busy waveform.
- 19. For information on port-to-port delay through RAM cells from writing port to reading port, refer to Read Timing with Busy waveform.
- 20. Test conditions used are Load 1.



Switching Characteristics Over the Operating Range^[13] (continued)

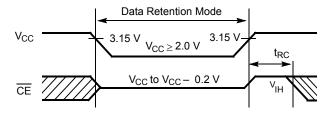
Parameter	Description	-	12	-15		Unit
		Min	Max	Min	Max	
Busy Timing ^[21]						
t _{BHC}	BUSY HIGH from CE HIGH	_	12	-	15	ns
t _{PS}	Port set-up for priority	5	-	5	-	ns
t _{WB}	R/W LOW after BUSY (Slave)	0	-	0	-	ns
t _{WH}	R/W HIGH after BUSY HIGH (Slave)	11	-	13	-	ns
t _{BDD} ^[22]	BUSY HIGH to data valid	_	12	-	15	ns
Interrupt Timing	[21]_					
t _{INS}	INT set time	_	12	-	15	ns
t _{INR}	INT reset time	_	12	-	15	ns
Semaphore Timi	ing					
t _{SOP}	SEM flag update pulse (OE or SEM)	10	_	10	-	ns
t _{SWRD}	SEM flag write to read time		_	5	-	ns
t _{SPS}	SEM flag contention window	5	_	5	_	ns
t _{SAA}	SEM address access time	_	12	_	15	ns

Data Retention Mode

The CY7C056V and CY7C057V are designed with battery backup in mind. Data retention voltage and supply current are guaranteed over temperature. The following rules ensure data retention:

- 1. Chip Enable (\overline{CE})^[23] must be held HIGH during data retention, within V_{DD} to V_{DD} 0.2 V.
- 2. $\overline{\text{CE}}$ must be kept between V_{DD} 0.2 V and 70% of V_{DD} during the power-up and power-down transitions.
- 3. The RAM can begin operation $>t_{RC}$ after V_{DD} reaches the minimum operating voltage (3.15 volts).

Timing



Parameter	Test Conditions ^[24]	Max	Unit
ICC _{DR1}	@ VDD _{DR} = 2 V	50	μΑ

- 21. Test conditions used are Load 1.

 22. t_{BDD} is a calculated parameter and is the greater of t_{WDD} – t_{PWE} (actual) or t_{DDD} – t_{SD} (actual).

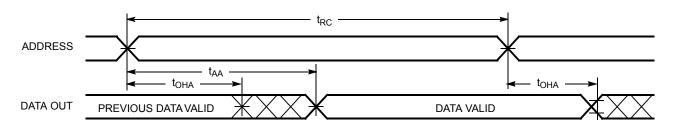
 23. \underline{CE} is LOW when $\overline{CE}_0 \leq V_{IL}$ and $\overline{CE}_1 \geq V_{IH}$.

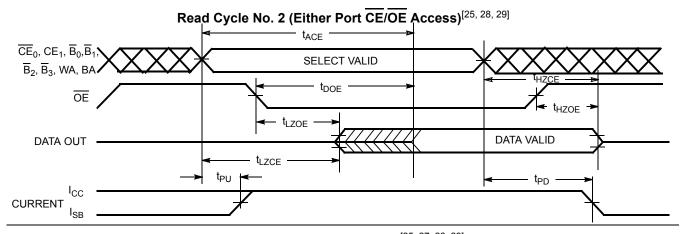
 24. $\overline{CE} = V_{DD}$, $V_{in} = V_{SS}$ to V_{DD} , $T_A = 25$ °C. This parameter is guaranteed but not tested.



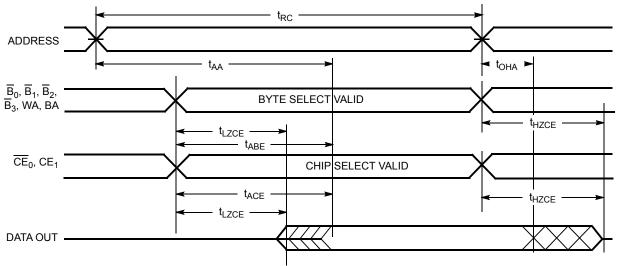
Switching Waveforms

Read Cycle No. 1 (Either Port Address Access)^[25, 26, 27]









- 25. R/W is HIGH for read cycles.
 26. <u>Device</u> is continuously selected. $\overline{CE}_0 = V_{IL}$, $\overline{CE}_1 = V_{IH}$, and \overline{B}_0 , \overline{B}_1 , \overline{B}_2 , \overline{B}_3 , WA, BA are valid. This waveform cannot be used for semaphore reads.

- 26. Device is continuously selected. SE₀ = S_L, S_L = N_L.

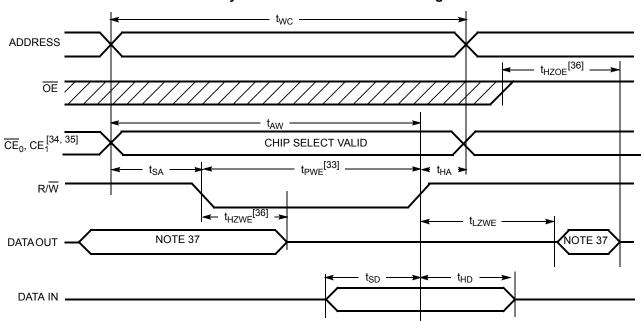
 27. OE = V_{IL}.

 28. Address valid prior to or coinciding with CE₀ transition LOW and CE₁ transition HIGH.

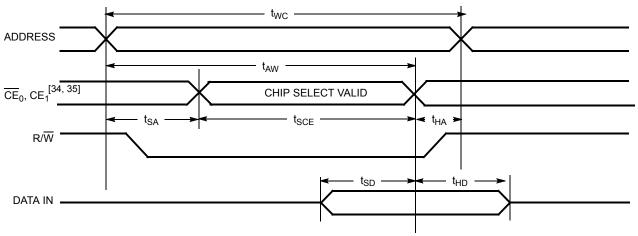
 29. To access RAM, CE₀ = V_{IL}, CE₄=V_{IL}, B₀, B₁, B₂, B₃, WA, BA are valid, and SEM = V_{IH}. To access semaphore, CE₀ = V_{IH}, CE₁=V_{IL} and SEM = V_{IL} or CE₀ and SEM=V_{IL}, and CE₁=B₀ = B₁ = B₂ = B₃, =V_{IH}.



Write Cycle No. 1: R/W Controlled Timing^[30, 31, 32, 33]



Write Cycle No. 2: CE Controlled Timing^[30, 31, 32, 38]



- 30. R/\overline{W} must be HIGH during all address transitions.
- 30. R/W must be HIGH during all address transitions.
 31. A write occurs during the overlap (t_{SCE} or t_{PWE}) of CE₀=V_{II} and CE₁=V_{IH} or SEM=V_{IL} and B₀₋₃ LOW.
 32. t_{Ha} is measured from the earlier of CE₀/CE₁ or R/W or (SEM or R/W) going HIGH at the end of Write Cycle.
 33. If OE is LOW during a R/W controlled write cycle, the write pulse width must be the larger of t_{PWE} or (t_{HZWE} + t_{SD}) to allow the I/O drivers to turn off and data to be placed on the bus for the required t_{SD}. If OE is HIGH during an R/W controlled write cycle, this requirement does not apply and the write pulse can be associated.

- to be placed on the bus for the required t_{SD}. If OE is HIGH during an R/W controlled write cycle, this requirement does not apply and the write pulse can be as short as the specified t_{PWE}.

 34. To access RAM, CE₀ = V_{IL}, CE₁=SEM = V_{IH}.

 35. To access byte B₀, CE₀ = V_{IL}, B₀ = V_{IL}, CE₁=SEM = V_{IH}.

 To access byte B₁, CE₀ = V_{IL}, B₁ = V_{IL}, CE₁=SEM = V_{IH}.

 To access byte B₂, CE₀ = V_{IL}, B₂ = V_{IL}, CE₁=SEM = V_{IH}.

 To access byte B₃, CE₀ = V_{IL}, B₃ = V_{IL}, CE₁=SEM = V_{IH}.

 To access byte B₃, CE₀ = V_{IL}, B₃ = V_{IL}, CE₁=SEM = V_{IH}.

 To access byte B₃, CE₀ = V_{IL}, B₃ = V_{IL}, CE₁=SEM = V_{IH}.

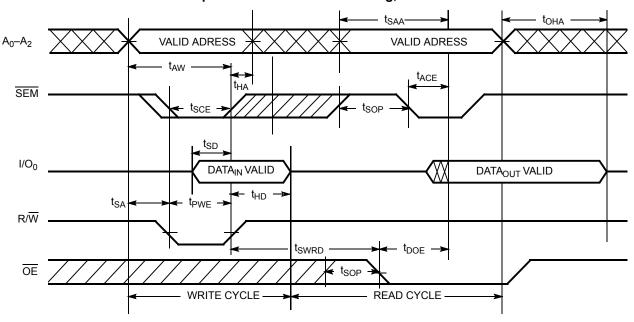
 36. Transition is measured ±150 mV from steady state with a 5-pF load (including scope and jig). This parameter is sampled and not 100% tested.

 37. During this period, the I/O pins are in the output state, and input signals must not be applied.

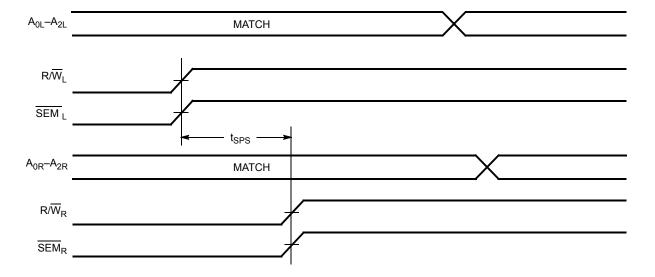
 38. If the CE₀ LOW and CE₁ HIGH or SEM LOW transition occurs simultaneously with or after the R/W LOW transition, the outputs remain in the high-impedance state.



Semaphore Read After Write Timing, Either Side^[39]



Timing Diagram of Semaphore Contention [40, 41, 42]

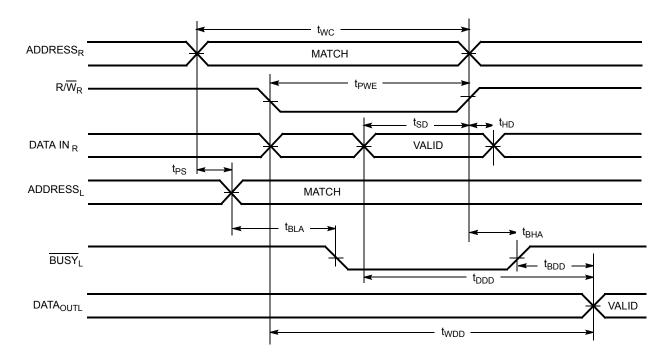


- 39. CE_0 = HIGH and CE_1 = LOW for the duration of the <u>ab</u>ove timing (both write and read cycle). 40. I/O_{0R} = I/O_{0L} = LOW (request semaphore); \overline{CE}_{0R} = \overline{CE}_{0L} = HIGH and CE_{1R} = CE_{1L} =LOW. 41. Semaphores are reset (available to both ports) at cycle start.

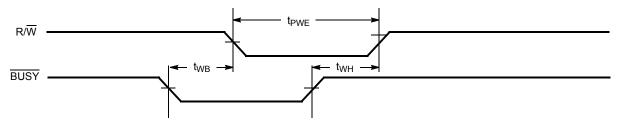
- 42. If t_{SPS} is violated, the semaphore will definitely be obtained by one side or the other, but which side will get the semaphore is unpredictable.



Timing Diagram of Write with BUSY $(M/S = HIGH)^{[43]}$

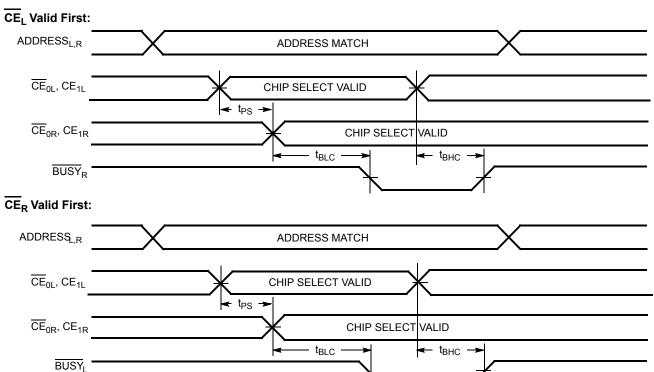


Write Timing with Busy Input (M/S = LOW)



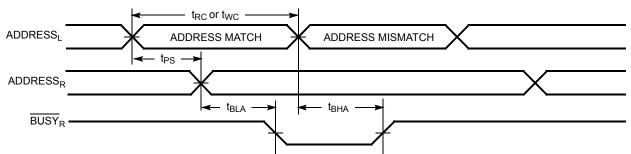


Busy Timing Diagram No. 1 (CE Arbitration)[44]

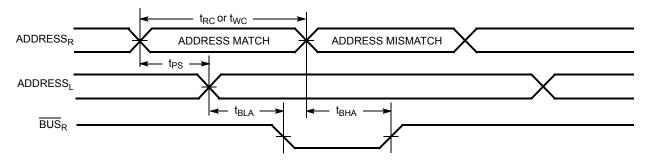


Busy Timing Diagram No. 2 (Address Arbitration)^[44]

Left Address Valid First:



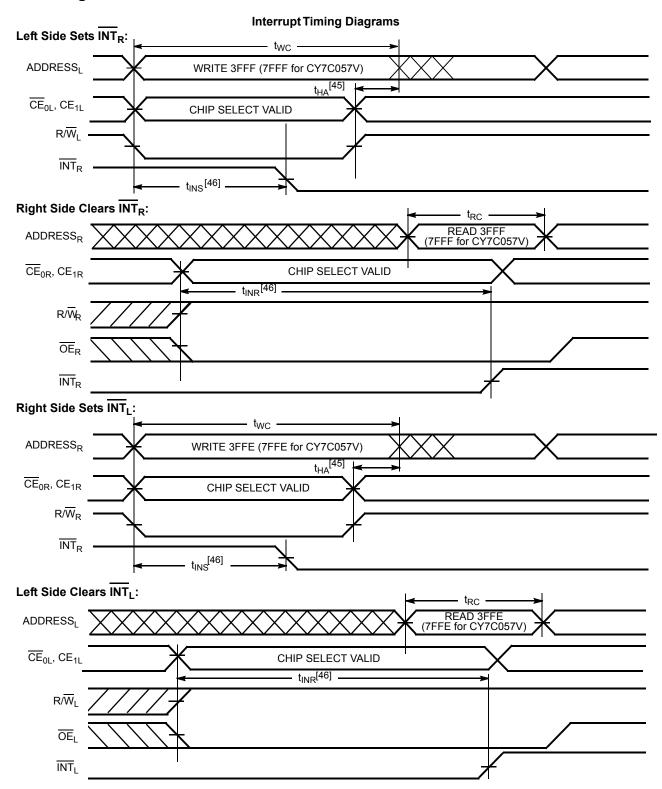
Right Address Valid First:



Note

^{44.} If t_{PS} is violated, the busy signal will be asserted on one side or the other, but there is no guarantee to which side BUSY will be asserted.





^{45.} t_{HA} depends on which enable pin $(\overline{CE}_{0L}/C\underline{E}_{1L})$ or $R_{1}\overline{W}_{L}$ is <u>deasserted</u> first. 46. t_{INS} or t_{INR} depends on which enable pin $(\overline{CE}_{0L}/CE_{1L})$ or $R_{1}\overline{W}_{L}$ is asserted last.



Architecture

The CY7C056V and CY7C057V consist of an array of 16K and 32K words of 36 bits each of dual-port RAM cells, I/O and address lines, and control signals ($\overline{\text{CE}_0}/\text{CE}_1$, $\overline{\text{OE}}$, R/W). These control pins permit independent access for reads or writes to any location in memory. To handle simultaneous writes/reads to the same location, a BUSY pin is provided on each port. Two Interrupt ($\overline{\text{INT}}$) pins can be utilized for port-to-port communication. Two Semaphore ($\overline{\text{SEM}}$) control pins are used for allocating shared resources. With the M/S pin, the devices can function as a master ($\overline{\text{BUSY}}$ pins are outputs) or as a slave ($\overline{\text{BUSY}}$ pins are inputs). The devices also have an automatic power-down feature controlled by $\overline{\text{CE}_0}/\text{CE}_1$. Each port is provided with its own Output Enable control ($\overline{\text{OE}}$), which allows data to be read from the device.

Functional Description

Write Operation

Data <u>m</u>ust be set up for a duration of t_{SD} before the rising edge of R/W in order to guarantee <u>a</u> valid write. A write operation is controlled by eith<u>er</u> the R/W pin (see Write Cycle No. 1 waveform) or the \overline{CE}_0 and CE_1 pins (see Write Cycle No. 2 waveform). Required inputs for non-contention operations are summarized in Table 1. If a location is being written to by one port and the opposite port attempts to read that location, a port-to-port flowthrough delay must occur before the data is read on the output; otherwise the data read is not deterministic. Data will be valid on the port t_{DDD} after the data is presented on the other port.

Read Operation

When reading the device, the user must assert both the \overline{OE} and $\overline{CE^{[3]}}$ pins. Data will be available t_{ACE} after \overline{CE} or t_{DOE} after \overline{OE} is asserted. If the user wishes to access a semaphore flag, then the SEM pin must be asserted instead of the $\overline{CE^{[3]}}$ pin, and \overline{OE} must also be asserted.

Interrupts

The upper two memory locations may be used for message passing. The highest memory location (3FFF for the CY7C056V, 7FFF for the CY7C057V) is the mailbox for the right port and the second-highest memory location (3FFE for the CY7C056V, 7FFE for the CY7C057V) is the mailbox for the left port. When one port writes to the other port's mailbox, an interrupt is generated to the owner. The interrupt is reset when the owner reads the contents of the mailbox. The message is user defined.

Each port can read the other port's mailbox without resetting the interrupt. The active state of the busy signal (to a port) prevents the port from setting the interrupt to the winning port. Also, an active busy to a port prevents that port from reading its own mailbox and, thus, resetting the interrupt to it.

If an application does not require message passing, do not connect the interrupt pin to the processor's interrupt request input pin. The operation of the interrupts and their interaction with Busy are summarized in Table 2.

Busy

The CY7C056V and CY7C057V provide on-chip arbitration to resolve simultaneous memory location access (contention). If

both ports' Chip Enables are asserted and an address match occurs within t_{PS} of each other, the busy logic will determine which port has access. If t_{PS} is violated, one port will definitely gain permission to the location, but it is not predictable which port will get that permission. \overline{BUSY} will be asserted t_{BLA} after an address match or t_{BLC} after \overline{CE} is taken LOW.

Master/Slave

A M/S pin is provided in order to expand the word width by configuring the device as either a master or a slave. The BUSY output of the master is connected to the BUSY input of the slave. This will allow the device to interface to a master device with no external components. Writing to slave devices must be delayed until after the BUSY input has settled (t_{BLC} or t_{BLA}), otherwise, the slave chip may begin a write cycle during a contention situation. When tied HIGH, the M/S pin allows the device to be used as a master and, therefore, the BUSY line is an output. BUSY can then be used to send the arbitration outcome to a slave.

Semaphore Operation

The CY7C056V and CY7C057V provide eight semaphore latches, which are separate from the dual-port memory locations. Semaphores are used to reserve resources that are shared between the two ports. The state of the semaphore indicates that a resource is in use. For example, if the left port wants to request a given resource, it sets a latch by writing a zero to a semaphore location. The left port then verifies its success in setting the latch by reading it. After writing to the semaphore, SEM or OE must be deasserted for t_{SOP} before attempting to read the semaphore. The semaphore value will be available t_{SWRD} + t_{DOE} after the rising edge of the semaphore write. If the left port was successful (reads a 0), it assumes control of the shared resource, otherwise (reads a 1) it assumes the right port has control and continues to poll the semaphore. When the right side has relinquished control of the semaphore (by writing a 1), the left side will succeed in gaining control of the semaphore. If the left side no longer requires the semaphore, a one is written to cancel its request.

Semaphores are accessed by asserting $\overline{\text{SEM}}$ LOW. The $\overline{\text{SEM}}$ pin functions as a chip select for the semaphore latches. For normal semaphore access, $\overline{\text{CE}}^{[3]}$ must remain HIGH during $\overline{\text{SEM}}$ LOW. A $\overline{\text{CE}}$ active semaphore access is also available. The semaphore may be accessed through the right port with $\overline{\text{CE}}_{0R}/\text{CE}_{1R}$ active by asserting the Bus Match Select (BM) pin LOW and asserting the Bus Size Select (SIZE) pin HIGH. The semaphore may be accessed through the left port with $\overline{\text{CE}}_{0L}/\text{CE}_{1L}$ active by asserting all $\overline{\text{B}}_{0-3}$ Byte Select pins HIGH. $\overline{\text{A}}_{0-2}$ represents the semaphore address. $\overline{\text{OE}}$ and $\overline{\text{R/W}}$ are used in the same manner as a normal memory access. When writing or reading a semaphore, the other address pins have no effect.

When writing to the semaphore, only I/O_0 is used. If a zero is written to the left port of an available semaphore, a 1 will appear at the same semaphore address on the right port. That semaphore can now only be modified by the port showing 0 (the left port in this case). If the left port now relinquishes control by writing a 1 to the semaphore, the semaphore will be set to 1 for both ports. However, if the right port had requested the semaphore (written a 0) while the left port had control, the right port would immediately own the semaphore as soon as the left port released it. Table 3 shows sample semaphore operations.



Table 1. Non-Contending Read/Write $^{[47]}$

Inputs				Outputs		
CE	R/W	OE	$\overline{B}_0, \overline{B}_1, \overline{B}_2, \overline{B}_3$	SEM	I/O ₀ –I/O ₃₅	Operation
Н	Х	Х	Х	Н	High Z	Deselected: Power-down
Х	Х	Х	All H	Н	High Z	Deselected: Power-down
L	L	Χ	H/L	Н	Data in and High Z	Write to selected bytes only
L	L	Х	All L	Н	Data in	Write to all bytes
L	Н	L	H/L	Н	Data out and High Z	Read selected bytes only
L	Н	L	All L	Н	Data out	Read all bytes
Х	Х	Н	Х	Х	High Z	Outputs disabled
Н	Н	L	Х	L	Data out	Read data in semaphore flag
Х	Н	L	All H	L	Data out	Read data in semaphore flag
Н		Х	Х	L	Data in	Write D _{IN0} into semaphore flag
Х		Х	All H	L	Data in	Write D _{IN0} into semaphore flag
L	Х	Χ	Any L	L		Not allowed

Table 2. Interrupt Operation Example (assumes $\overline{\text{BUSY}}_{\text{L}} = \overline{\text{BUSY}}_{\text{R}} = \text{HIGH})^{[47, 48]}$

	Left Port				Left Port Right Port					
Function	R/W _L	CEL	OEL	A _{0L-13L}	INTL	R/W _R	CER	OE _R	A _{0R-13R}	INT _R
Set right INT _R flag	L	L	Х	3FFF	Х	Х	Х	Х	Х	L ^[50]
Reset right INT _R flag	Х	Х	Х	Х	Х	Х	L	L	3FFF	H ^[49]
Set left INT _L flag	Х	Х	Х	Х	L ^[49]	L	L	Х	3FFE	Х
Reset left INT _L flag	Х	L	L	3FFE	H ^[50]	Х	Х	Х	X	Х

Table 3. Semaphore Operation Example

Function	I/O ₀ -I/O ₈ Left	I/O ₀ -I/O ₈ Right	Status
No action	1	1	Semaphore free
Left port writes 0 to Semaphore	0	1	Left port has semaphore token
Right port writes 0 to Semaphore	0	1	No change. Right side has no write access to Semaphore
Left port writes 1 to Semaphore	1	0	Right port obtains semaphore token
Left port writes 0 to Semaphore	1	0	No change. Left port has no write access to semaphore
Right port writes 1 to Semaphore	0	1	Left port obtains semaphore token
Left port writes 1 to Semaphore	1	1	Semaphore free
Right port writes 0 to Semaphore	1	0	Right port has semaphore token
Right port writes 1 to Semaphore	1	1	Semaphore free
Left port writes 0 to Semaphore	0	1	Left port has semaphore token
Left port writes 1 to Semaphore	1	1	Semaphore free

 $[\]label{eq:Notes} \begin{array}{c} \textbf{Notes} \\ \textbf{47. CE} \text{ is LOW when } \overline{\textbf{CE}}_0 \leq \textbf{V}_{IL} \text{ and } \textbf{CE}_1 \geq \textbf{V}_{IH}. \\ \textbf{48. A}_{0L=14L} \text{ and } \textbf{A}_{0R-14R}, \textbf{7FFF/7FFE} \text{ for the CY7C057V}. \\ \textbf{49. If } \underline{\textbf{BUSY}}_R = \textbf{L}, \text{ then no change}. \\ \textbf{50. If } \underline{\textbf{BUSY}}_L = \textbf{L}, \text{ then no change}. \\ \end{array}$



Right Port Configuration[51, 52, 53]

ВМ	SIZE	Configuration	I/O Pins Used
0	0	x36 (standard)	I/O ₀₋₃₅
0	1	x36 (CE active SEM mode)	I/O ₀₋₃₅
1	0	x18	I/O ₀₋₁₇
1	1	х9	I/O ₀₋₈

Right Port Operation

Configuration	WA	ВА	Data Accessed ^[54]	I/O Pins Used
x36	Х	Х	DQ ₀₋₃₅	I/O ₀₋₃₅
x18	0	Х	DQ ₀₋₁₇	I/O ₀₋₁₇
x18	1	Х	DQ ₁₈₋₃₅	I/O ₀₋₁₇
х9	0	0	DQ ₀₋₈	I/O ₀₋₈
х9	0	1	DQ ₉₋₁₇	I/O ₀₋₈
х9	1	0	DQ ₁₈₋₂₆	I/O ₀₋₈
х9	1	1	DQ ₂₇₋₃₅	I/O ₀₋₈

Left Port Operation

Control Pin	Effect
BO	I/O ₀₋₈ Byte control
B1	I/O ₉₋₁₇ Byte control
B2	I/O _{18–26} Byte control
B3	I/O _{27–35} Byte control

When reading a semaphore, data lines 0 through 8 output the semaphore value. The read value is latched in an output register to prevent the semaphore from changing state during a write from the other port. If both ports attempt to access the semaphore within t_{SPS} of each other, the semaphore will definitely be obtained by one side or the other, but there is no guarantee which side will control the semaphore.

When reading a semaphore, data lines 0 through 8 output the semaphore value. The read value is latched in an output register to prevent the semaphore from changing state during a write from the other port. If both ports attempt to access the semaphore within $t_{\rm SPS}$ of each other, the semaphore will definitely be obtained by one side or the other, but there is no guarantee which side will control the semaphore.

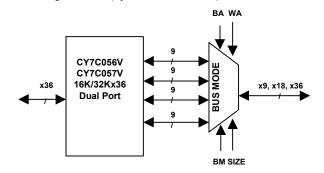
When reading a semaphore, data lines 0 through 8 output the semaphore value. The read value is latched in an output register to prevent the semaphore from changing state during a write from the other port. If both ports attempt to access the semaphore within $t_{\rm SPS}$ of each other, the semaphore will definitely be obtained by one side or the other, but there is no guarantee which side will control the semaphore.

Notes

- 51. BM and SIZE must be configured one clock cycle before operation is guaranteed.
- 52. In x36 mode WA and BA pins are "Don't Care."
- 53. In x18 mode BA pin is a "Don't Care."
- 54. DQ represents data output of the chip.

Bus Match Operation

The right port of the CY7C057V 32Kx36 dual-port SRAM can be configured in a 36-bit long-word, 18-bit word, or 9-bit byte format for data I/O. The data lines are divided into four lanes, each consisting of 9 bits (byte-size data lines).



The bus match select (BM) pin works with bus size select (SIZE) to select bus width (long-word, word, or byte) for the right port of the dual-port device. The data sequencing arrangement is selected using the word address (WA) and byte address (BA) input pins. A logic "0" applied to both the bus match select (BM)



pin and to the bus size select (SIZE) pin will select long-word (36-bit) operation. A logic "1" level applied to the bus match select (BM) pin will enable either byte or word bus width operation on the right port I/Os depending on the logic level applied to the SIZE pin. The level of bus match select (BM) must be static throughout device operation.

Normally, the bus size select (SIZE) pin would have no standard-cycle application when BM = LOW and the device is in long-word (36-bit) operation. A "special" mode has been added however to disable ALL right port I/Os while the chip is active. This I/O disable mode is implemented when SIZE is forced to a logic "1" while BM is at a logic "0". It allows the bus-matched port to support a chip enable "Don't care" semaphore read/write access similar to that provided on the left port of the device when all Byte Select (\overline{B}_{0-3}) control inputs are deselected.

The bus size select (SIZE) pin selects either a byte or word data arrangement on the right port when the bus match select (BM) pin is HIGH. A logic "1" on the SIZE pin when the BM pin is HIGH selects a byte bus (9-bit) data arrangement). A logic "0" on the SIZE pin when the BM pin is HIGH selects a word bus (18-bit) data arrangement. The level of the bus size select (SIZE) must also be static throughout normal device operation.

Long-Word (36-bit) Operation

Bus match select (BM) and bus size select (SIZE) set to a logic "0" will enable standard cycle long-word (36-bit) operation. In this mode, the right port's I/O operates essentially in an identical fashion as does the left port of the dual-port SRAM. However no byte select control is available. All 36 bits of the long-word are shifted into and out of the right port's I/O buffer stages. All read and write timing parameters may be identical with respect to the two data ports. When the right port is configured for a long-word size, word address (WA), and byte Address (BA) pins have no application and their inputs are "Don't Care" [55] for the external user.

Word (18-bit) Operation

Word (18-bit) bus sizing operation is enabled when bus match select (BM) is set to a logic "1" and the bus sze select (SIZE) pin is set to a logic "0." In this mode, 18 bits of data are ported through I/O_{0R-17R} . The level applied to the word address (WA) pin during word bus size operation determines whether the most-significant or least-significant data bits are ported through the I/O_{0R-17R} pins in an Upper word/Lower word select fashion (note that when the right port is configured for word size operation, the Byte Address pin has no application and its input is "Don't care" [55]).

Device operation is accomplished by treating the WA pin as an additional address input and using standard cycle address and data setup/hold times. When transferring data in word (18-bit) bus match format, the unused $I/O_{18R-35R}$ pins are three-stated.

Byte (9-bit) Operation

Byte (9-bit) bus sizing operation is enabled when bus match select (BM) is set to a logic "1" and the bus size select (SIZE) pin is set to a logic "1." In this mode, data is ported through I/O_{0R-8R} in four groups of 9-bit bytes. A particular 9-bit byte group is selected according to the levels applied to the word address (WA) and byte address (BA) input pins.

I/Os	Rank	WA	BA
I/O _{27R-35R}	Upper-MSB	1	1
I/O _{18R-26R}	Lower-MSB	1	0
I/O _{9R-17R}	Upper-MSB	0	1
I/O _{0R-8R}	Lower-MSB	0	0

Device operation is accomplished by treating the word address (WA) pin and the byte address (BA) pins as additional address inputs having standard cycle address and data set-up/hold times. When transferring data in byte (9-bit) bus match format, the unused I/O_{9R-35R} pins are three-stated.

Note

^{555.} Even though a logic level applied to a "Don't Care" input will not change the logical operation of the dual-port, inputs that are temporarily a "Don't Care" (along with unused inputs) must not be allowed to float. They must be forced either HIGH or LOW.



Ordering Information

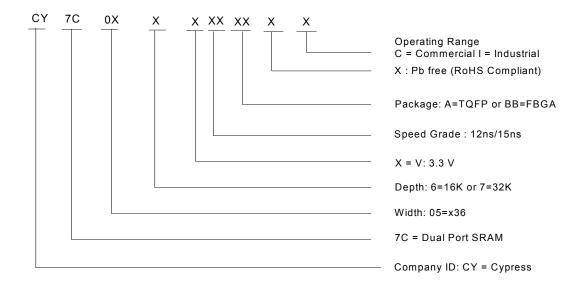
16K x 36 3.3 V Asynchronous Dual Port SRAM

Speed (ns)	Ordering Code	Package Name	Package Type	Operating Range
15	CY7C056V-15AXC	A144	144-Pin Pb-free Thin Quad Flat Pack	Commercial

32K x 36 3.3 V Asynchronous Dual Port SRAM

Speed (ns)	Ordering Code	Package Name	Package Type	Operating Range
12	CY7C057V-12AC	A144	144-Pin Thin Quad Flat Pack	Commercial
	CY7C057V-12AXC	A144	144-Pin Pb-free Thin Quad Flat Pack	
15	CY7C057V-15AXC	A144	144-Pin Pb-free Thin Quad Flat Pack	Commercial
	CY7C057V-15AXI	A144	144-Pin Pb-free Thin Quad Flat Pack	Industrial
	CY7C057V-15BBI	BB172	172-Ball Ball Grid Array (BGA)	Industrial
	CY7C057V-15BBXC	BB172	172-Ball Ball Grid Array (BGA)	

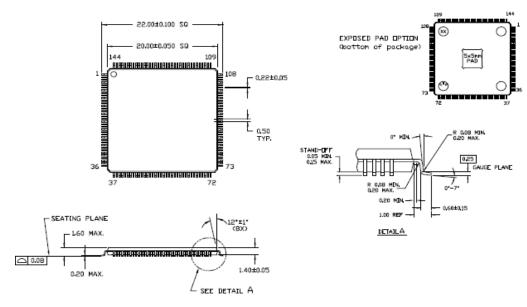
Ordering Code Definition





Package Diagrams

Figure 3. 144-Pin Plastic Thin Quad Flat Pack (TQFP) A144



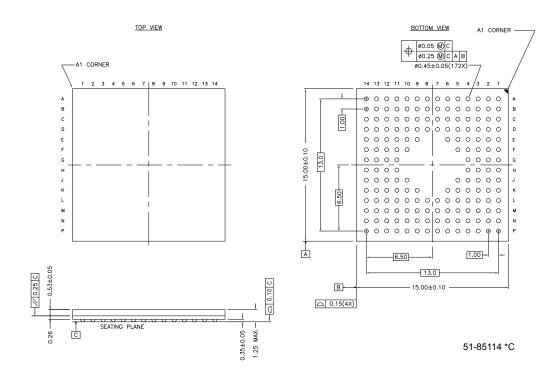
(DIMENSIONS ARE IN MILLIMETERS)

51-85047 *C



Package Diagrams (continued)

Figure 4. 172-Ball FBGA (15 x 15 x 1.25 mm) BB172



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Acronyms

Acronym	Description
CMOS	complementary metal oxide semiconductor
TQFP	thin quad plastic flatpack
I/O	input/output
SRAM	static random access memory
BGA	ball grid array

Document Conventions

Units of Measure

Symbol	Unit of Measure
ns	nano seconds
V	Volts
μA	micro Amperes
mA	milli Amperes
Ω	Ohms
mV	milli Volts
MHz	Mega Hertz
pF	pico Farad
W	Watts
°C	degree Celcius



Document History Page

REV.	ECN NO.	Issue Date	Orig. of Change	Description of Change
**	110214	12/16/01	SZV	Change from Spec number: 38-00742 to 38-06055
*A	122305	12/27/02	RBI	Power up requirements added to Maximum Ratings Information
*B	393770	See ECN	YIM	Added Pb-Free Logo Added Pb-Free parts to ordering information: CY7C056V-12AXC, CY7C056V-15AXC, CY7C057V-12AXC, CY7C057V-15AXC, CY7CO57V-15AXI
*C	2897217	03/22/2010	RAME	Updated Ordering Information Updated Package Diagrams
*D	3093365	11/25/2010	ADMU	Removed part CY7C057V-15BBC Added part CY7C057V-15AXI Updated datasheet as per new template Added Acronyms and Units of Measure table Added Ordering Code Definition Updated all footnotes.
*E	3210221	03/30/2011	ADMU	Removed parts CY7C056V-15AC and CY7C057V-12BBC from the Ordering Information table.



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Page 26 of 26